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# Determination of the turning point of cache efficiency in computer networks with logic $E\tau$

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# Abstract

Object caches are used to minimize data traffic in various areas of Information Technology, including computer networks. In this scenario, they are usually hosted in proxies, storing page objects (texts, figures, among others), and implementing access control policies. Its correct operation can provide a significant gain of performance in data exchange since it allows the immediate response of requested resources. This study aims to discuss the different states of the efficiency of two distinct types of computer network caches and determine the changing dynamics of these states with Logic  $E\tau$ .

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Keywords: Type your keywords here, separated by semicolons ;

# 1. Introduction

The task to analyze the efficiency of a computer network cache will be based on Logic  $E\tau$ , composed of memory and disk cache. The hit rates for each cache type will be evaluated over a given period and will formulate hypotheses

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about any changes in its state.

Computer networks can be considered one of the main ways of transmitting information, if not the main one. Monitoring the information has become a critical factor in technology sectors [1] and is part of the backbone of Information Technology in various educational institutions and companies. A computer network consists of several connected hosts, which can be represented by a desktop, a laptop, a smartphone, wearable devices, biomedical sensors, among others [2][3]. In such a heterogeneous client environment, dynamic content adaptation and delivery services are becoming an essential requirement for the new Internet service infrastructure [4].

This paper is organized as follows: in Section 2, the methodology and a comparison of the evaluated project with existing anomaly detection models are discussed; in Section 3, the evaluated network attributes used are presented; in Section 4, the basic concepts of Logic  $E\tau$  are introduced; section 5 discusses the development of the project considering the network attributes and Logic  $E\tau$ ; section 6 presents the results and discussion, and the conclusion is found in Section 7.

#### 2. Methodology and comparison with existing models

In order to compare with existing models, it is necessary to present the materials and tools used in this project. The information was acquired from a proxy server, representing a gateway for external networks. The evaluated network is used by the academic staff of an educational institution, exclusively by its students and teachers.

The network attributes were normalized, and concepts of the Logic  $E\tau$  were applied, considering favorable and contrary evidence for each scenario.

Unlike [5] and [6], which rely on simulated data to emulate a real network environment and the synthetic generation of anomalies, this project employs continuously gathered data from the operation of a real working network for the learning process. Another difference from [7] is that this project uses Digital Signature of Network Segment using Flow Analysis (DSNSF), which establishes a profile for the normal behavior of a network segment by considering the history of its movement. A possible problem: when a real-time system is not considered for this kind of task, any changes in the network layout or its availability may impact the analyzer learning process, since the history may not represent the actual state of the network.

Another significant difference from the work of [8], in which data traffic was used as an analytical measure, without distinction of individual attributes that could represent different operating situations of the network, is that, in this project, the network attributes are treated individually.

## 3. Evaluated network attributes

As mentioned by [9]: "The amount of information that travels across the Internet has increased dramatically over the past decades because of the huge growth in the number of internet users." Also, the growing number of services and applications, as well as the many advances in Information Technology, make networks and information systems essential for the survival of all educational enterprises, organizations, and institutions [10]. The growth of the global computer network also leads to an increase in the complexity of its infrastructure. Thus, the classical methods of network analysis may not be the adequate ones for this scenario[11].

Computer networks use routers to communicate with each other. As mentioned by [12], "Routing is the process of sending data packets from the host of origin to the destination host, which is performed by the routers." Some elements may be interesting for network traffic analysis. The following analyzable attributes were considered:

- TCP\_MEM\_HIT/200
- TCP\_HIT/200

The first attribute (TCP\_MEM\_HIT/200) represents the requests that had a positive response in the search for objects stored in memory. The second attribute (TCP\_HIT/200) represents the response for objects stored in the disk.

# 4. The Logic Eτ

There are high levels of uncertainty when monitoring data in computer networks. The argument for this assertion is based on the principle that user actions are presented as random elements [13]. Therefore, the use of a nonclassical logic becomes an option. Logic  $E\tau$  can be a viable technique to search for indications of problems during the regular operation of the network or by intentional elements [14] [15]. In the latter case, the problems may be caused by misuse or malicious software [16].

According to [17]: "The atomic formulas of the Logic E $\tau$  are the type  $p(\mu, \lambda)$ , where  $(\mu, \lambda) \in [0, 1]^2$  ([0, 1] is the real unit interval) and p denote a propositional variable". Therefore, among several readings,  $p(\mu, \lambda)$  can be intuitively read: "It is assumed that the favorable evidence of p is  $\mu$  and the contrary evidence of p is  $\lambda$ ." Thus, we have, for instance, the following particular readings:

- $p_{(1.0, 0.0)}$  can be read as a true proposition
- $p_{(0.0, 1.0)}$  as false
- $p_{(1.0, 1.0)}$  as inconsistent
- $p_{(0.0, 0.0)}$  as paracomplete, and
- $p_{(0.5, 0.5)}$  as an indefinite proposition

The uncertainty and certainty degrees associated with  $(\mu, \lambda)$  are defined [18][19]:

- Uncertainty Degree:  $G_{un}(\mu, \lambda) = \mu + \lambda 1 \ (0 \le \mu, \lambda \le 1);$
- Certainty Degree:  $G_{ce}(\mu, \lambda) = \mu \lambda \ (0 \le \mu, \lambda \le 1);$

An order relation is defined on  $[0, 1]^2$ :  $(\mu_1, \lambda_1) \leq (\mu_2 \lambda_2) \Leftrightarrow \mu_1 \leq \mu_2$  and  $\lambda_2 \leq \lambda_1$ , forming a lattice, which is symbolized by  $\tau$ .

With the degrees of certainty and uncertainty, one can determine the following 12 output states, shown in Table 1:

Extreme states	Symbol	Non-extreme states	Symbol
True	V	Quasi-true tending to Inconsistent	QV→T
False	F	Quasi-true tending to Paracomplete	$QV \rightarrow \perp$
Inconsistent	Т	Quasi-false tending to Inconsistent	QF→T
Paracomplete	$\perp$	Quasi-false tending to Paracomplete	QF→⊥
<u> </u>	•	Quasi-inconsistent tending to True	QT→V
		Quasi-inconsistent tending to False	QT→F
		Quasi-paracomplete tending to True	Q⊥→V
		Quasi-paracomplete tending to False	Q⊥→F

Table 1. Extreme and the Non-Extreme S	tates
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Extreme and non-extreme states are shown in Fig. 1:



Figure 1: Extreme and non-extreme states of the Lattice  $\boldsymbol{\tau}$ 

In Fig 2, the states, along with certainty and uncertainty degrees, are shown, as well as the control values.



Figure 2: Certainty / Uncertainty degrees with decision states of the Lattice

# 5. Development

The cache performance analysis has considered an entire period of local network operation, from 8 AM to 10 PM. For this analysis, there are no influences from the cache created on the previous day. Each operation cycle is reset, i.e., they are deleted for the creation of a new one. This operation is performed automatically at the end of the previous day's operation.

The research was based on data obtained in a computer network of a public educational institution, composed by approximately 230 hosts, and it covers not only the students and professors of this institution (academic area) but also the employees (administrative area). Therefore, they can be described as sectors with different demands and behaviors.

The fundamental element in this analysis is represented by the accesses made by the network users, that is, the situations in which the users have searched for some resource in the network. Usually, it is difficult to predict the actions of the users, who usually present a random behavior pattern in the network. Therefore, the cache structuring presents an even greater element of uncertainty, which must be considered.

The number of accesses throughout the day can be represented by the "Total Requisition" field, which is counted at every hour of operation of a computer network. The access values can be represented by Fig. 3:



Figure 3: "Total Requisition" values throughout the day operation

Similarly, a representation of the cache hits, considering the data of memory and disk, can be represented by figure 4:



#### Figure 4: Memory and disk cache hits

The cache obtained its best performance starting at 15:00, presenting a tendency of continuous growth until 18:00 when it reaches its peak performance (Figure 4). Then, it suffers a sudden performance drop, remaining at low levels from 19:00 to 22:00. On the other hand, we can verify, in Figure 3, that the "Total Requisition" field remained stable starting at 15:00, with little variation compared to the cache, hits performance.

#### 6. Results and Discussion

From the values normalized above, the Logic  $E\tau$  was applied to verify in which state this abrupt variation could be framed. As favorable evidence, the values of hits in-memory cache (TCP\_MEM\_HIT/200) were applied. As contrary evidence, there was the application of the disk cache hits (TCP\_HIT/200). It is notorious that access to resources in the memory is significantly faster than on disk, which allows a more effective response to the requests. The Lattice  $\tau$  was generated by presenting the various analyzed cache hits, represented by Fig. 5:



Figure 5: Favourable and contrary evidence from cache hits on Lattice  $\tau$ 

From the presented states in the lattice described in Figure 5, the only two points in which the cache operation can be considered efficient are represented by the intervals of 17:00 and 18:00. Next, there is a sudden loss of performance starting at 19:00, when both favorable and contrary evidence decline.

In comparison with the attribute "Total Requisitions," it is observed that there is no increase, at any moment, in traffic data that could justify this significant change. Instead, data decrease is observed in the network.

Therefore, after achieving its best performance, both disk and memory caches were no longer able to handle the requests of stored objects. Such a phenomenon may be explained by the fact that, after 18:00, there may have been a significant change in the object profile requested by the users. Considering that it consists of an academic network, a change of context must be evaluated, considering that the interval comprises the entrance of the night shift students.

### 7. Conclusion

With this study, it has been possible not only to determine the maximum performance in absolute values of the network cache but also to quantify its respective components.

Also, it has been possible to determine that the leading cause of the performance loss of the cache was much more due to the use of the network by the students, than in comparison to the institution employees, who presented relatively constant behavior. By the time students start the night shift, the cache has its performance decreased quickly, and it does not have enough time to recover its past performance.

Also, through the individual analysis of the memory and disk cache, it has been possible to make a qualitative analysis, since the speed of the first is significantly higher than the second's, which is a desirable scenario. It is understood that the memory cache attending the user's requests should represent a more responsive and efficient service.

For the network manager, the determination of when the cache loses its efficiency is fundamental, since it can lead to actions to correct this loss of performance as soon as possible. For example, a possible solution for the analyzed scenario is the separation of the student and employee data traffic, since the latter is undetermined by the less predictable behavior of the first. A simple way to achieve this goal without modifying the physical network layout is the use of VLANs, a feature available in the vast majority of layer two switches on the market.

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